

Generic search algorithms

Tree search can be used if the state space is a tree, otherwise **graph search** must be used. All search algorithms (BFS, DFS, uniform-cost, A*, etc) are variations of one of these (usually graph search). The only difference between tree search and graph search is that tree search does not need to store the explored set, because we are guaranteed never to attempt to visit the same state twice.

TREE-SEARCH(problem):

Initialize the frontier using the initial state of problem

Loop forever:

If frontier is empty, return failure

Choose a leaf node from the frontier and remove it (from the frontier)

If the node contains a goal state, return the corresponding solution

If the node did not contain a goal state, expand the chosen node, adding the resulting nodes to the frontier

GRAPH-SEARCH(problem):

Initialize the frontier using the initial state of the problem

Initialize the explored set to empty

Loop forever:

If frontier is empty, return failure

Choose a leaf node from the frontier and remove it (from the frontier)

If the node contains a goal state, return the corresponding solution

Add the node to the explored set

If the node did not contain a goal state, expand the chosen node, adding the resulting nodes to the frontier, but only if the resulting node is not already in the frontier or the explored set

Uninformed search algorithms

Breadth-first search: Run the generic graph search algorithm with the frontier stored as a (LIFO) queue.

Depth-first search: Run the generic graph search algorithm with the frontier stored as a (FIFO) stack.

Uniform cost search

We have a node data structure that contains a state, a path-cost (also known as g), a pointer to the parent node, and the action that generated this state from the parent state.

UNIFORM-COST-SEARCH(problem) // aka Dijkstra's algorithm

node \leftarrow a new node corresponding to the initial state, with path-cost (g) = 0

frontier \leftarrow a priority queue of nodes sorted by path-cost (g), initialized to contain only node

explored \leftarrow an empty set of states

loop forever:

if frontier is empty, return failure

node \leftarrow pop(frontier) // remove lowest path-cost node from frontier (smallest g)

if IS-GOAL(node.state), then return the corresponding solution

add node.state to the explored set

for each action in ACTIONS(node.state):

child_node \leftarrow new node with child_node.state = RESULT(node.state, action)

child_node.g = node.g + COST(node.state, action, child_node.state)

child_node.action = action

child_node.parent = node

if child_node.state is not in explored or frontier:

add child_node to frontier

else if child_node.state is already in frontier, but child_node.g is smaller than the frontier's:

replace that frontier node with child_node