

Dijkstra's Algorithm (book page 655)

```
void dijkstra(Graph g, Vertex start, Vertex finish)
{
    create min-priority queue Q

    for each vertex v in the graph:
        dist[v] = infinity
        prev[v] = undefined

    dist[start] = 0
    Q.insert(start, 0)

    while Q is not empty:
        u = Q.extract_minimum()           # We now "visit" vertex u.

        if u == finish: break

        for each neighbor v of u:
            alt = dist[u] + weight(u, v)
            if alt < dist[v]
                dist[v] = alt
                prev[v] = u
                if Q.contains(v)
                    Q.decrease_priority(v, alt)
                else
                    Q.insert(v, alt)

    Final path length is dist[finish].

    Traverse prev[] array starting from prev[finish] in reverse order back
    to start vertex to get final path from start to finish.
}
```

Note: during the for each neighbor v of u step, the algorithm will reconsider nodes it has already visited before (thereby opening the possibility of a cycle). However, for a situation like this, $\text{dist}[u] + \text{weight}(u, v)$ will always be bigger than $\text{dist}[v]$, so the cycles will be ignored anyway. However, some Dijkstra's Algorithm implementations explicitly keep track of which vertices have been visited already and modify the for each neighbor v of u step to skip over any vertex v that has already been visited earlier.