CS342: Bioinformatics Assembling a Genome



- 1. <u>Overlap</u> Build the overlap graph
- <u>Layout</u> Bundle stretches of the overlap graph into *contigs*.
- **3.** <u>Consensus</u> Pick the most likely nucleotide sequence for each contig.



Sometimes additional information can be used to begin to *scaffold* or order together contigs.

Overlap graph is big and messy!!! Contigs don't just pop out.

Part of the overlap graph for:

To_every_thing_turn_turn_turn_there_is_a_season

Fragment Length: 7 Minimum Overlap: 4

You can remove **redundant** information (edges) from the graph.



The layout step can also deal with problematic subgraphs, likely the result of errors.



Mismatch could be due to sequencing error or repeat. Since the path through **b** ends abruptly we might conclude it's an error and prune **b**.

Contigs correspond to non-branching stretches in the simplified graph





At each position, ask: what nucleotide (and/or gap) is here?

Complications: (a) sequencing error, (b) ploidy

Say the true genotype is AG, but we have a high sequencing error rate and only about 6 reads covering the position.

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De Bruijn Graph Assembly

A formulation conceptually similar to overlapping/SCS, but has some potentially helpful properties not shared by SCS.

The graph of a sequence

	GACGG	ì								
	ACGG	iC								
	CGG	iCG								
	GG	iCGG								
	G	icggc								
		CGGCG								
		GGCGC		- 1	Гhе	complete	set	of	16	5-mers
		GCGC	Α							
		CGC	AC							
		GC	ACG							
		C	ACG	G						
Now we can construct a graph where:		4	ACG	GC						
•	Each 5-mer is a node		CG G	GCG GCG(2					
•	There is a directed edge from a k -mer that shares its $(k - 1)$ -base suffix with (k - 1)-base prefix of another k -mer	the	G	GCGC. CGC	A CAA					

A read-overlap graph



The problem is *How to infer the original sequence from this graph?*

The rules of our game

- Every node, *k*-mer, can be used exactly once
- The object is to find a path along edges that visits every node (vertex) one time
- This game was invented in the mid 1800's by a mathematician called Sir William Hamilton
 A version of Hamilton's game:







Finding a Hamiltonian Path in a graph



Example Problem: Hamiltonian Path Approach

S = { ATG AGG TGC TCC GTC GGT GCA CAG }

ATG AGG TGC TCC GTC GGT GCA CAG



ATGCAGGTCC

Path visited every VERTEX once

More Complicated Example: Hamiltonian Path Approach

A more complicated graph:



More Complicated Example: Hamiltonian Path Approach *S* = { ATG TGG TGC GTG GGC GCA GCG CGT }



Another way to represent our *k*-mers in a graph

- Rather than making each k-mer a node, let's try making them an edge
- That seems odd, but it is related to the overlap idea



- Think of the k-mer as the edge connecting a prefix to a suffix
- This leads to a series of simple graphs
- Then combine all nodes with the same label

A De Bruijn Graph

This rather odd graph is called the "De Bruijn" graph; was named after a famous mathematician.

The problem is *How to infer the original sequence from this graph?*



The rules of our new game

- Every edge, *k*-mer, can be used exactly once
- The object is to find a path in the graph that uses each edge only one time
- This game was invented in the late 1700's by a mathematician called Leonhard Euler



Leonhard Euler

A version of Euler's game:



Bridges of Königsberg Find a city tour that crosses every bridge just once

Example Problem: Eulerian Path Approach

S = { ATG, TGG, TGC, GTG, GGC, GCA, GCG, CGT }

Vertices correspond to (k - 1) – mers : { AT, TG, GC, GG, GT, CA, CG }

Edges correspond to *k* – mers from *S*



Example Problem: Eulerian Path Approach

S = { AT, TG, GC, GG, GT, CA, CG } corresponds to two different paths:

